

Load Optimization with Coverage and Connectivity for Wireless Sensor Networks

Shreyas R¹, Abdul Haq N²

^{1,2}Dept. of ECE Engineering, REVA University, Karnataka, Bangalore-India, 560064

Abstract - The main important function of wireless sensor networks is coverage, sensing and connectivity. Life time of the network is mainly reduced due to the coverage problem imposed by the densely congested sensors in the wireless sensor networks (WSNs). To overcome such issues, in this paper the proposed novel maximum connected load balancing cover tree (MCLCT) to get entire coverage and base station connectivity of every node by formulating load balanced routing cover trees. For the limited energy of sensor nodes, we are gathering these nodes in to subgroups, which is efficient of monitoring all discrete point of interest (DPOIs). And it is a maximum cover tree problem, which is non-probabilistic complete problem. The MCLCT consists of two strategies: a coverage optimizing recursive heuristic for coverage and the probabilistic load balancing for the determination of the routing path. By using the MCLCT less amount of energy consumption can be done by balancing the load evenly on the nodes. The simulation results show that the energy efficiency and the connectivity of the nodes.

Key Words: (Size 10 & Bold) Key word1, Key word2, Key word3, etc (Minimum 5 to 8 key words)...

1. INTRODUCTION

Fast growing [10] and demanding in the field in today's world of computing is data gathering. Sensors give a easy solution to these applications particularly in the remote and low-maintenance areas where conservative approaches prove to be very expensive. Networking of these unattended sensors is expected to have significant impact on the effectiveness of many military and civil applications, such as combat field observation, security and disaster management. Power is a very scarce reserve for such sensor systems. Even in WSN power is considered as scarce resource and should be efficiently used.

The main aim of this paper to provide the better coverage and connectivity [1] of the WSNs. The deployment of the wireless sensor nodes near the discrete point of interest (DPOIs), the region where there is no efficient signal or remote area. And how each DPOI will be covered is the main issue. This is one of the problem in WSNs. And scheduling the process is the main feature which will determine the time of activation [1] and time of inactivation of the nodes. By using the localization technologies, it is possible to locate the sensor nodes. Grouping the sensor nodes in to a maximum no. of cover sets which are disjoint sets or non-disjoint sets is proved to be NP-complete problem. This

nodes [1] in the cover sets can monitor all the DPOIs. To achieve the longer life time alternate activation of this cover set will be done.

The connected set cover (CSC) [2] which targets to find a maximum number of cover sets. The sensor nodes will be activated only when it is connected to the base station and it is a non-probabilistic complete problem. To overcome this issues they have come up with two algorithms: an integer programming based heuristic and a breadth heuristic. And to deal with connected cover set problem distributed heuristic was also implemented which is on minimum spanning tree.

Sleep scheduling mechanisms and coverage mechanisms are used to extend the life time of the network. Virtual robust spanning tree and modified virtual robust spanning tree is used for connectivity preservation [3]. This is best when compared to cardei method because it reduces energy consumption than the cardei method. By using VRST and MVRST has less depth and data latency. And uses existing point coverage methods to reduce energy consumption of sensor node and to construct a load balanced tree to fight against network failure.

And some solutions [4] to forming k-connected coverage of targets with minimum number of active nodes. This is a k-connected argumentation problem which is a target coverage problem, it is a NP-hard. To overcome this issues, heuristic algorithm is used. This aforementioned [2-3-4] studies tells about the CSC problem. They are only practical about sensing and communications. And it will not effect to the distance between the nodes. And equal amount of energy is consumed for both sensing and relaying. Based on the size of the data to be transmitted only energy consumption should be done.

Connected target coverage problem [5] is similar to maximum cover. Algorithm known as communication weighted greedy cover (CWGC) is used. Optimized connected coverage heuristic (OCCH) is one of the algorithm which is based on general coverage method. Generate the connected cover sets, where the critical nodes would not serve as relaying nodes by assigning different weights to edges between nodes. By doing so, [5] the major energy of the critical nodes could be conserved so as to prolong the network lifetime. The experimental result showed that the network lifetime acquired by the OCCH is longer than those of the aforementioned approaches, including CWGC [6], Greedy-CSC [2], and GIECC [2].

Large no. of sensors are deployed for the surveillance applications to get [15] efficient coverage. Due to the unexpected node failures, they focus on the fault tolerance and energy efficiency in sensor coverage. To overcome the node failures, the proposed distributed technique to facilitate switching between active covers without the urge for node synchronization. This [15] framework is called as “resilient online coverage” framework. It allows only sensor node to take control over the degree of redundancy and surveillance in its region according to the current network condition.

2. DESIGN EQUATIONS AND ANALYSIS

This section [1] tells about the main optimization problem is maximum cover tree problem. The load of sensor nodes are formulated and providing load balance among the nodes. Let n number of sensor nodes are there in WSN which is denoted by S, Where $S=\{S1,S2,S3.....Sn\}$ and m number of DPOIs is denoted by P. And $P=\{P1,P2,P3,.....Pm\}$. Assume at least one sensor node will cover one dpoi. And initial energy is denoted by E0. And let τ denotes the operational time of the cover tree. The nodes[1] undergoes three modes of operations sensing, relaying and sleep mode/inactive mode. And $ts(si, \tau)$, $tr(si, \tau)$, and $tinact(si, \tau)$ will also denote the operational period of time of the cover tree.

$$ts(si, \tau) + tr(si, \tau) + tinact(si, \tau) = \tau \quad (1)$$

For the energy consumption for communication and sensing operations both free space and multipath fading channels are used depending on the distance between the transmitter and receiver. So the energy $e_t^{i,j}$ consumed for transferring information depends on the transmitter and receiver. The free space model is used when the distance is less than the threshold, or otherwise the multipath model is used.

$$e_t^{i,j} = \begin{cases} ect + ecfs \cdot d_{ij}^2, & dij < d_0 \\ ect + ecmp \cdot d_{ij}^4, & dij \geq d_0 \end{cases} \quad (2)$$

‘ect’ is represented as consumption of energy by transmitted circuit per data bit. ‘ecfs’ and ‘ecmp’ denote the energy consumption for the power amplifier per data bit for free space and multipath model respectively. Euclidian distance is denoted by dij between the node Si and Sj. d0 denoted as,

$$d_0 = \sqrt{\frac{ecfs}{ecmp}} \quad (3)$$

2.1 MAXIMUM CONNECTED LOAD BALANCING COVER TREE ALGORITHM

The main agenda of the MCT problem is to build several connected cover trees. [10] By doing so, a longer network lifetime and full coverage can be acquired. The MCT problem

is a complicated NP-Complete problem, so finding a suboptimal solution is a generic approach in order to decrease the time of computation. The proposed MCLCT is composed of two sub-strategies: a coverage-optimizing recursive (COR) heuristic and a probabilistic load-balancing (PLB) strategy. The COR heuristic aims at finding a maximum number of disjoint sets of nodes, which can be achieved by one of the sensor nodes (such as the sink node). In each disjoint set, the nodes are able to monitor all the DPOIs together. That is, the COR heuristic focuses on dealing with the full coverage preservation issue. Moreover, the PLB strategy is used to figure out the appropriate path from each node to the BS after the disjoint sets are initiated.

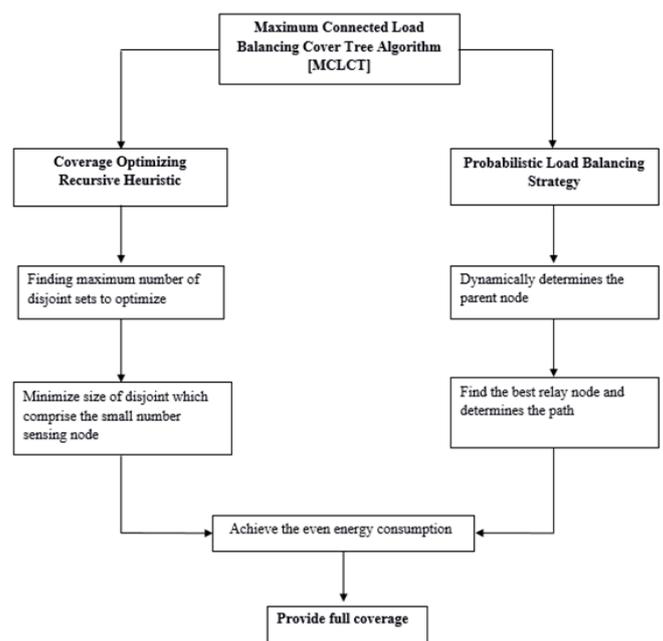


Fig.1 Flowchart of methodology used in MCLCT

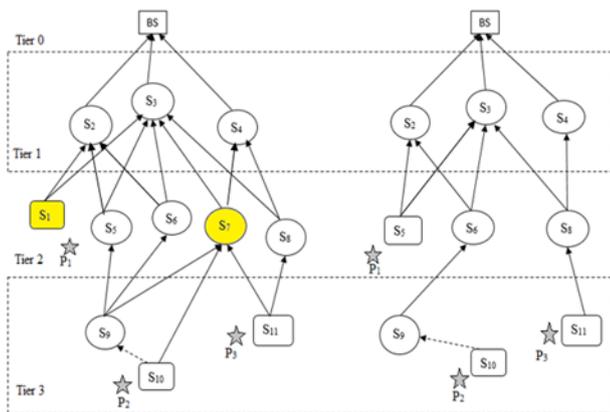
The main aim[1] of this paper to provide full coverage and connectivity in the wireless sensor networks and which is comprises of:

Lost coverage and base station connectivity of sensing nodes can be dynamically recovered because some of the factors such as weather change, man-made interference, congestion.. etc., some set nodes are reserved which is called as residual nodes which are used recover the lost data by coverage optimizing recursive heuristic strategy.

Load balance among the wireless sensor nodes, when the nodes will try to choose the best candidate parent nodes by using probabilistic load balancing strategy, they will also help to build the overall load-balanced dynamic cover trees. And so the network hot spot and congestion can be avoided, and a longer network lifetime can be increased. At the same tier maintaining connections between the neighbor nodes.

Size of every disjoint set is minimized by the existing maximum load balancing cover tree selects as few nodes as possible and put them into the disjoint sets using the coverage optimizing recursive heuristic. Only necessary sensing nodes will be activated during the operation in order to reduce energy waste.

And by avoiding the simultaneous operations playing as sensing and relaying by the nodes, a node plays one role at one time as sensing or relaying.



$$S_{res} = \{S_2, S_3, S_4, S_5, S_6, S_7, S_8, S_9\}$$

(a) $T(\tau)$ at τ'

$$S_{res} = \{S_2, S_3, S_4, S_5, S_6, S_8, S_9\}$$

(b) $T(\tau)$ at $\tau' + 1$

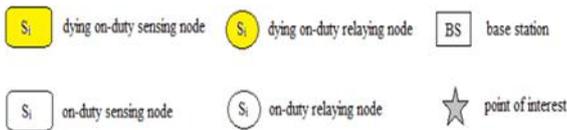


Fig.2 Coverage recovery by MCLCT

Fig. 2(a). Tells about the [1]dynamic coverage recovery in which P1, P2, P3 are the DPOIs and which there are 3 sensing nodes that will monitor this DPOIs. 8 relaying nodes used to transfer the sensed data. And s1, s10, and s11 are the sensing nodes that form the disjoint set S2, S3, S4, S5, S6, S7, S8, S9 form the set Srsd.) At the $T(\tau)$ at τ' there will be connection established between s9 and s10. And s10 has a sole candidate parent node (s7) in the lower tier

(Tier-2). Under such circumstances, maximum connected load balancing cover tree algorithm will conduct a dynamic coverage recovery, because the p1 will not be covered, which is due to the malfunction of S1. Fig. 2 (b) tells about the new topology of $T(\tau)$ at $\tau' + 1$. As seen in the fig.2(b) node S5 originally serving as a relaying node would be changed to serve as a sensing node at $\tau' + 1$, because its coverage resembles with that of S1. And without connecting to S10 at

$\tau' + 1$, the sensed data of S10 will not be transferred to the base station. And by the dynamic compensation mechanism of the maximum connected load balancing cover tree for coverage and connectivity, the roles of sensor nodes can be efficiently determined.

The main objective of the probabilistic load balancing strategy is to fulfill the load balance among the wireless sensor nodes by sharing load equally among the nodes. The main function of the sensing nodes to monitor the DPOIs and transfer the data that sensed by it. By undergoing this, there will be less amount of energy consumption can be maintained. And equation (4), the residual energy is represented as $ef(\tau')$ of the node S_i at $\tau'.$ α denotes exponential factor used to set the response strength of the curve from energy to distance ratio. The sensing node s_i to the candidate parent node, the forwarding probability can be expressed as,

$$\tilde{p}(S_i, S_{rj}, \tau') = \frac{\left(\frac{e_{rj}(\tau')}{d_{i,rj}}\right)^\alpha}{\sum_{K \in Prt(S_i)} \left(\frac{e_k(\tau')}{d_{i,k}}\right)^\alpha} \quad (4)$$

Equation (4), is further simplified. Before performing the forwarding probability for the relaying node, the normalization of the residual energy of the candidate parent nodes using their expected loads. It is done with the expected load, the weight of the one-hop transmission path from the relaying node S_k to its candidate parent node $[srh \in Prt(sk)$ at τ'].

$$\tilde{p}(S_k, S_{rj}, \tau') = \frac{Dev(weight(sk, srj, \tau')) - 2.minDev}{\sum_{Srheprt(s_k)} (Dev(weight(sk, srj, \tau')) - 2.minDev)} \quad (5)$$

And by using equation (5), the forwarding probability that S_k tries to relay data to the candidate parent node S_{rh} with bigger positive weight deviation will be high. By doing so, forwarding probability [1] that S_k attempts to transfer data to S_{rh} with a larger negative weight deviation will be too smaller. Therefore, tendency of load balance can be found when S_k chooses the best one from the candidate parent nodes to relay data, by using the weight deviations. Relaying on serving as sensing node or a relaying node, the forwarding probabilities for candidate parent nodes can be computed according to the equation (4) or (5). By this idea of probability assignment is suitable for wireless sensor networks. Without depending on every sensor node instead of on a sink node or a BS with a higher computation capability and unlimited resource.

3. RESULTS AND COMPARISON

In this section, the simulation results regarding the performance of the proposed MCLCT and compare it with previously mentioned algorithms which include CWGC [1], OCCH [2], GIECC [4]. All the simulations are carried out on the network simulation -2 platform.

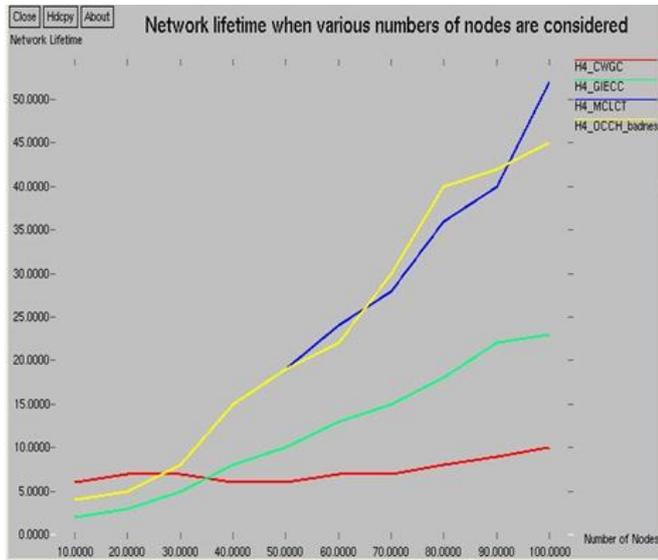


Fig 3. Network lifetime when various numbers of nodes are considered (10 to 100 nodes).

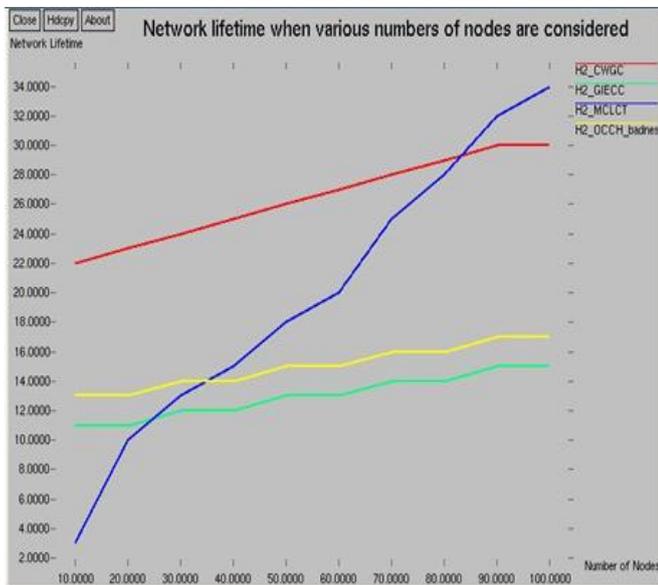


Fig 4. Network lifetime when various numbers of nodes are considered

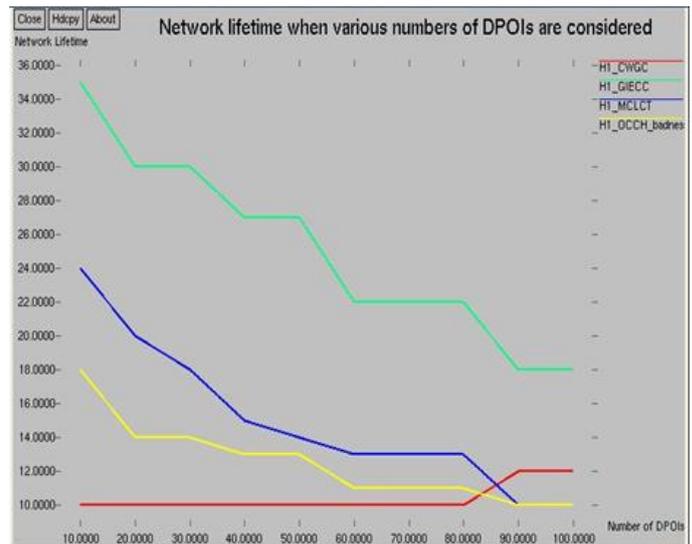


Fig 5. Network lifetime when various number of DPOIs are considered.

4. CONCLUSION

The simulation results show that network life duration is efficiently increased because of the traffic, the burden generated by the sensing nodes is dynamically organized to various relaying nodes. Further, for power conservation, the MCLCT will definitely performs well-arranged power utilization for nodes through the construction of dynamic cover tree.

REFERENCES

- [1] [1]. Chia-Pang Chen, Member, IEEE, S.C. Mukhopadhyay, Fellow, IEEE, Cheng-Long Chuang, Member, IEEE, Maw-Yang Liu, Member, IEEE, and Joe-Air Jiang, Senior Member, IEEE “Efficient Coverage and Connectivity Preservation with Load Balance for Wireless Sensor Networks”
- [2] [2] I. Cardei and M. Cardei “Energy-efficient connected – coverage in wireless sensor networks,” Int. J. Sens. Netw., vol. 3, no. 3, pp. 201–210, 2008..
- [3] [3]. P. Ostovari, M. Dehghan, and J. Wu, “Connected point coverage in wireless sensor networks using robust spanning trees,” in Proc. 31st International Conference on Distributed Computing Systems Workshops (ICDSC’11), 2011, pp. 287–293.
- [4] [4]. N. Jaggi and A. A. Abouzeid, “Energy-efficient connected coverage in wireless sensor networks,” in Proc. 4th Asian International Mobile Computing Conference (AMOC’06), 2006, pp. 77–86.
- [5] [5]. D. Zorbas and C. Douligeris, “Connected coverage in WSNs based on critical targets,” Comput. Netw., vol. 55, pp. 1412–1425, 2011.
- [6] [6]. Q. Zhao and M. Gurusamy, “Lifetime maximization for connected target coverage in wireless sensor networks,” IEEE/ACM Trans. Netw., vol. 16, no. 6, pp. 1378–1391, 2008.

- [7] [7]. X. Han, X.G. Cao, E. L. Loyd, and C.-C. Shen, "Fault-tolerant relay node placement in heterogeneous wireless sensor networks," *IEEE T. Mobile Comput.*, vol. 9, no. 5, pp. 643–656, 2010.
- [8] [8]. A. Krause, R. Rajagopal, A. Gupta, and C. Guestrin, "Simultaneous optimization of sensor placements and balanced schedules," *IEEE Trans. Autom. Control*, vol. 56, no. 10, pp. 2390–2405, 2011.
- [9] [9]. D. Yang, S. Misra, X. Fang, G. Xue, and J. Zhang, "Two-tiered constrained relay node placement in wireless sensor networks: computational complexity and efficient approximations," *IEEE Trans. Mobile Comput.*, vol. 11, no. 8, pp. 1399–1411, 2012.
- [10] [10]. P.Gowtham¹, P.Vivek Karthick², "Optimized Coverage and Efficient Load Balancing Algorithm for WSNs-A Survey [11].H. M. Ammari and S. K. Das, "Centralized and clustered k-coverage protocols for wireless sensor networks," *IEEE T. Comput.*, vol. 61, no. 1, pp. 118–133, 2012.
- [11] [12]. M. Ashouri, Z. Zali, S.R. Mousavi, and M.R. Hashemi, "New optimal solution to disjoint set K-coverage for lifetime extension in wireless sensor networks," *IET Wirel. Sens. Syst.*, vol. 2, no. 1, pp. 31–39, 2012
- [12] [13]. Y. Lin, J. Zhang, H.S.-H. Chung, W. H. Ip, Y. Li, and Y.-H. Shi, "An ant colony optimization approach for maximizing the lifetime of heterogeneous wireless sensor networks," *IEEE Trans. Syst., Man, Cybern. C, Appl. ,Rev.*, vol. 42, no. 3, pp. 408–420, 2012.
- [13] [14]. X. He, H. Yanh, and X. Gui, "The maximum coverage set calculated algorithm for WSN area coverage," *J. Netw.*, vol. 5, no. 6, pp. 650–657, 2010.
- [14] [15]. O. M. Younis, M. M. Krunz, and S. Ramasubramanian, "ROC: resilient online coverage for surveillance applications," *IEEE/ACM Trans. Netw.*, vol. 19, no. 1, pp.251–264, 2011.