

A Distributed Time Triggered Control for a Feedback Control System

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Abstract - The paper presents a distributed time triggered control using time triggered CAN (TTCAN) for a closed-loop control system. The balance of computing load is the rule to separate the control system into functional parts for controllers. The controllers are operated by time triggered control with TTCAN communication. The challenge is minimizing the response time to regulate the controlled system. It is done by adjusting the processing cycle that depends on the tasks' processing time, data transmission time and the clock synchronizing drift. The system operation is managed by the communication and operation schedules which are also able to deal with several aperiodic events. The designs are experimented on an inverted rotary pendulum with the processing period of three milliseconds. The tasks and communication activities are triggered at the predefined time points and processed within the designed intervals. That makes the pendulum working stably at upright position.

Key Words: distributed control, time triggered, feedback control system, Inverted rotary pendulum, TTCAN

1. INTRODUCTION

Distributed control is applied for a distributed network system which consists of multiple subsystems [1]. The subsystems called nodes receive data and transmit the output via communication network to implement the tasks. To regulate a closed-loop control system, real time control is required. It means the tasks in the system must be completed before the deadlines [2].

Time triggered CAN has wide application in industrial application with advantages such as high dependability, deterministic communication and low latency jitters of transmitting messages [3-4]. In the protocol, all events such as transmitting or receiving in the system are activated by time segment elapsing [3]. Despite of drawbacks such as lack on flexibility and the restrictive design process, it is significantly suitable for a closed-loop control system that works autonomously and has little interaction with human [3].

In the paper, a distributed time triggered control is designed for an inverted rotary pendulum (IRP) using TTCAN as communication mean. The challenge is minimizing the response time to make the calculated value from the control program still valuable to regulate the IRP. To shorten the response time, there are two problems to solve. The first one is all critical tasks that directly affect to the control process such as data collection, data processing and actuator control

must be implemented inside predefined intervals. The second one is the communication has to perform deterministically with little latency. Moreover, the distributed control can handle the events such as adjusting the input from the operator, or displaying the IRP's states.

The distributed control is performed on a set of controllers which cannot individually process all the tasks. The tasks are assigned to controllers based on the balance of computation load. To manage the system operation, communication and operation schedules are built. And the controllers' clocks are synchronized.

FreeRTOS is deployed for the controllers as operation system. It gives the priority for the task which should perform ahead if two tasks perform at the same time [5]. Therefore, ensuring the critical tasks are implemented at predefined time points.

The paper's content is as follow: section 2 discusses about distributed control, section 3 mentions about Time Triggered Control. Section 4 describes the experiment and results, and the conclusion is discussed in section 5.

2. DISTRIBUTED CONTROL

To design a distributed control, the controlled system is analyzed the functionality then separated into tasks. Based on the rule of balance of computing load, the tasks are grouped and assigned to the controllers.

2.1 Inverted Rotary Pendulum

IRP is a two links system with two rotational joints as shown on Fig -1. It consists of an arm and a bar that rotate around the Z- and Y-axis with two angles α and ϕ , respectively. In steady state, mass M stands at the upright position and the arm stays at a desired position.

The variations of α and ϕ (the IRP's outputs) are detected by two optical encoders. From the errors between the angles and desired inputs, the controllers calculate amount of energy used and direction for the DC motor. In a limited time interval the data's processing and communication must be completed to response to the angles' variations, thus keep the mass M at the upright position. Moreover, a Human - Machine Interface (HMI) is required to interact with operator.

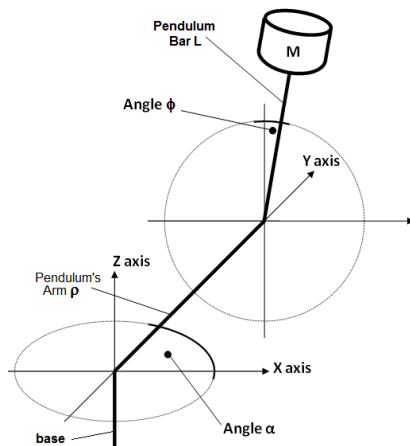


Fig -1: IRP's structure

Based on the above analysis, the IRP control program consists of the main parts: Sensors signal collecting, data processing, DC – motor controlling, and Human – Machine Interface.

2.2 Distributed Control

From the previous analysis, the control program is separated into individual tasks, and the order of the data needed to control the IRP. The tasks are continued to analyze further following.

The first task is to handle the sensor's signals. The sensors detect any change of the IRP's outputs, and adjust the *output variables*. The variable's values then are transmitted to other nodes at the time points which are predefined by the communication schedule. The task causes the hosted controller working under highly frequent interrupt due to the sensor's signals.

The second task is to process the data including output variables and the reference input. To calculate the PWM values for the DC motor, the task contains many parameters such as control coefficients, IRP's parameters and states [6]. Therefore, it has high computing load in the IRP control program and is dedicated a controller to perform.

The third task is to control the DC motor. Similarity to the first task, the controller which hosts the task implements other tasks under highly frequent interrupt. All three tasks above are required hard real-time performance [2] to keep the IRP operating stable.

Another function of the IRP is to create a human – machine interface. It contains two parts: one is for displaying the IRP's states on a display device, and another one is for handling the operator command via an input device such as a push bottom matrix. The latter task deals with emergency situation or adjusting the input, thus it is required hard real-time control. There is no consequence if violating the deadline, thus the former task classified as soft real-time control [2].

The order of tasks' implementation is: the first task receives the sensors' signals, update the values to the output

variables then send the values to other controllers. The second task gets the data from the first task to calculate the PWM value. The third one then uses the PWM value to export corresponding energy level to the DC motor.

From the analysis, the distributed control for the IRP is designed with four controllers and shown on the Fig -2. The second task has highest computing load, thus is assigned to one controller – the third one. The first and the third tasks have a common feature that is invoked with high frequency. Therefore, both are assigned for the first and four controllers, respectively. The HMI function is in charge of the second controllers.

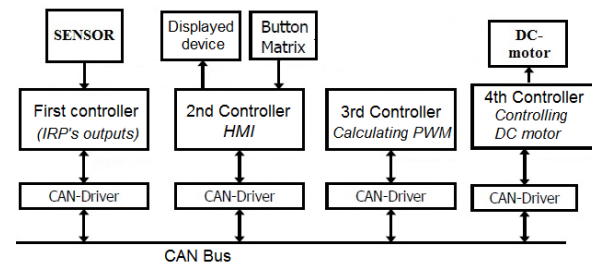


Fig -2: Distributed Control

The messages are exchanged in the network by TTCAN communication which is discussed in the next section.

3. TIME TRIGGERED CONTROL

Deploying the distributed control, two kinds of schedules are set up to manage the communication and operation of the controllers. The first one called *Communication Schedule* is used to manage TTCAN communication. The second one is *Operation Schedule*, which controls the controllers' operation.

3.1 Time Triggered CAN

In the protocol, communication activities are triggered by the time elapsing and happened periodically. The period is called *basic cycle*, which consist of multiple *time slots* [3]. Each time slot is started by a Time Mark. The communication schedule is designed to activate a particular event at the right Time Mark and point out the related controllers. Time slot interval is then analyzed following.

- *Communication Schedule*

It manages all activities on the communication channel in one basic cycle. Time slots can vary in size, however, should be enough for message frame transmission [2-4]. They are set equal in this case. From the distributed control, there are three kinds of message frames needed: *Sensor*, *Command*, and *PWM*. Sensor frame stores the system output values – the sensors' values. Command frame contains the reference input and command of operator. PWM frame stores the PWM value, which controls amount of energy for the DC motor. To synchronize the controllers' clocks, it is needed a message frame called reference frame – *MasTim* [3]. This frame is managed by the third controller, which is now called Time

Master. Therefore, there are four message frames which are corresponding to four time slots. One more time slot is for redundancy and another one is to distinguish two consecutive basic cycles. All the time slots and the activities in four nodes are shown on Fig -3.

At Time Mark '0', the starting point of a basic cycle, *TimMas* frame is sent from Time Master, which manages the global time of the system. After receiving, all the non-master nodes use the frame's data to synchronize the global time [3].

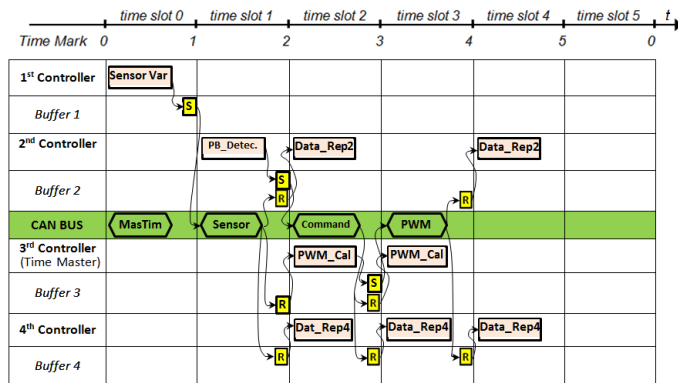


Fig -3: Communication Schedule

In the same time slot, Sensor frame which stores the current IRP's outputs is prepared to transmit. In time slot '1', it is sent to other controllers. The frame's data is shown on the display device, update working ranges of the pendulum, and used to calculate the PWM value. As the calculation is done, its result is transmitted via PWM frame in time slot '3'. The fourth controller uses this data to export amount of energy to regulate the motor. If any operator's command is detected, it is sent to the third and fourth controller via Command frame in time slots '2'. It will stop the motor immediately if any incident occurs. The time slots '4' is redundancy for emergency situation and future use.

• Time Slot Interval T_I

Due to the data order of controlling the IRP, the time slot interval is required long enough for both frame transmission and tasks' processing. Therefore, two cases are considered to calculate the interval. The longer one will be used to set up the time slot interval. Moreover, the clock synchronizing time drifts T_d and redundancy time for retransmission are also taken into account.

Case 1: Based on Transmission Time

The interval is required to complete a transmission of the longest frame. Due to the only difference in the data field, Sensor frame is the longest frame with six bytes with the transmitted time called T_s . Combining with the retransmission redundancy time (equal to T_s) and the time drift of synchronizing the clocks T_c , the minimum time slot interval T_I is:

$$T_I = 2T_s + T_c \tag{1}$$

Case 2: Based on Processing Time

The processing time is considered for *PWM_Calculation* task which has highest computing load. The task is performed in the third controller - Time Master, thus there is no time drift. Therefore, the time slot interval is greater than processing time T_p :

$$T_I \geq T_p \tag{2}$$

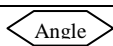
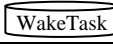
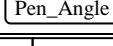
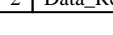
Based on the set up of CAN bus and the computing power of the controllers, T_I calculated in the equation (1) and (2) will be different. The greater one is chosen for the time slot interval.

3.2. Controller Operation Schedules

This section describes how to organize the tasks in the controllers to meet two requirements: the critical tasks with hard real-time operation must be completed before the deadline and they must be synchronized with *Communication Schedule*. The synchronization is managed by the task *Operation Schedule* which is invoked at each Time Mark; then it triggers other tasks to work. The activated method is to send a signal to predefined queues that in turn waking up the specified tasks [4]. Due to the role, Operation Schedule has highest priority in the controllers.

To analyze the controllers, Table -1 presents symbols and the meaning used to describe tasks' structure.

Table -1: Symbol used for tasks' structure description

Symbol	Description
	Message frame
	Queue or semaphore
	Variable in controller
	Task (priority task name) or Interrupt Service Routine (ISR)

• First Controller

There are two tasks: *Operation Schedule 1* and *Data_Preparation*. The data called variables are the IRP's outputs – two angles *Pen_A* and *Arm_A*. They are updated by Interrupt Service Routines (ISR) which detect the signals' change from the encoders. At time mark '0', *Operation Schedule 1* sends a signal to the queue *Que* to invoke the task *Data_Preparation* whose duties is to prepares the data for *Sensor* frame. The frame is transmitted in time slot 1. In the time slot '0', the controller receives *MasTim* frame from Time Master to synchronize the clock. The task's structure is shown in Fig -4.

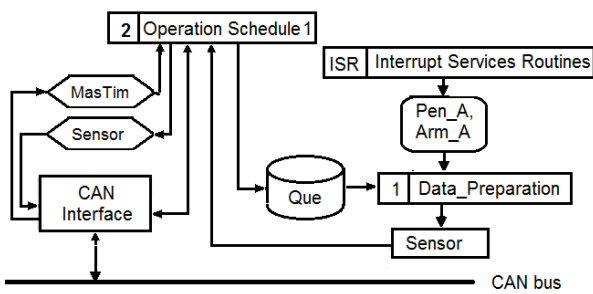


Fig -4: Tasks' structure in the first controller

• Second Controller

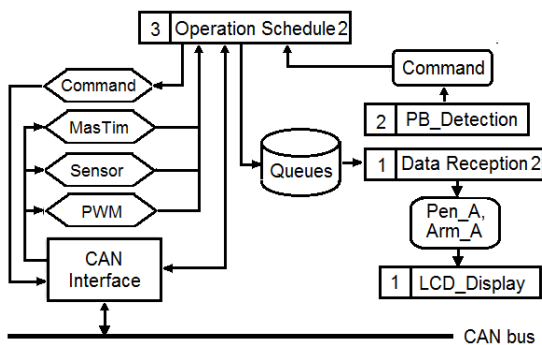


Fig -5: Task's structure in the second controller

Due to managing the Human – Machine Interface, the controller is connected with two peripheral devices: a LCD and Pushing Button Matrix. To display IRP's states, it needs two tasks: *Data Reception 2* and *LCD Display*. The former one is waken up at time mark 1, 2, and 4 by *Operation Schedule 2* to handles all the incoming frames (MasTim, Sensor, PWM). The frames' data is then updated for the displayed variables. The latter one manages the LCD and shows the information of the screen. The third task *PB_Detection* is used to sense the Button Matrix, recognize the operator's command, and pack it into a Command frame. The frame is transmitted to the destination at time mark '2'. The task's composition is shown in Fig -5.

• Third Controller

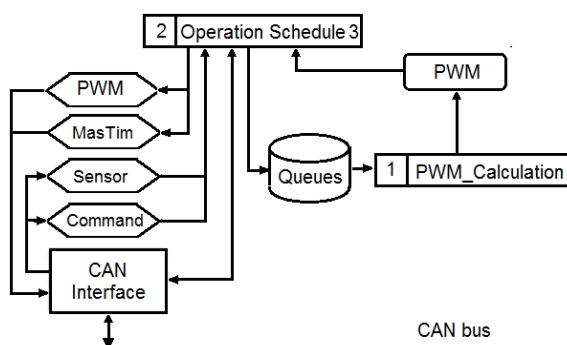


Fig -6: Task's structure in the third controller

It is Time Master and used to calculate PMW value. As *Command* and *Sensor* frames have arrived, *PWM_Calculation* is waken up by *Operation Schedule 3* to process the data. The result is packed in *PWM* frame and sent to the fourth

controller. *Operation Schedule 3* task sends MasTim frame at time slot '0'.

• Fourth Controller

The controller's duty is to control amount of engery and rotated direction for the DC-motor. In order to control two working phases of the IRP, there are three tasks (*Data Reception*, *Upright Controller*, *SwingUp Controller*) cooperating shown on Fig -7.

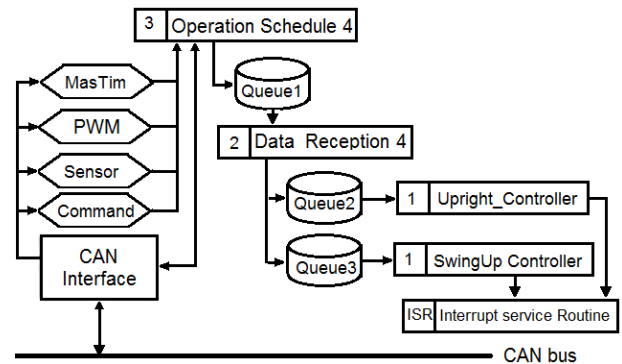


Fig -7: Task's structure of the fourth controller

As message frames come, *Operation Schedule 4* sends a signal to *Queue1* to wake up task *Data Reception 4*. Based on the received data (*PWM*, *Sensor*, *Command*), it processes and sends another signal to *Queue2* or *Queue3* to activate *Upright_Controller* task or *SwingUp_Controller* task, respectively. The tasks then affect to *Interrupt Service Routine* to control the energy level.

The distributed time triggered control is then applied to an IRP model. The experiement is implemented to determine the appropriate control parameters.

4. EXPERIMENT AND RESULTS

The design is applied on a IRP model, which was made in the university's workshop. In the experiment, the signals of CAN bus and the states of the critical tasks are measured by MS07014B Mixed Signal Oscilloscope. The results are discussed afterward.

4.1 Experiment

The IRP model is shown on Fig -8 which is driven by a DC motor 12V. The dynamic energy is transmitted via a gearbox, a belt transmission with the transmitted ratios 18 and 4, respectively. The revolution of the encoders' is 2000 values/rev. The pendulum has mass 0.03 kg. The arm and pendulum's bar have the lengths 0.13m and 0.4m. And the inertia moment of the base on which the arm is fixed is 0.002 Kgm².

The model is controlled by a circuit board on Fig -9 which contains a LCD, a push botton matrix and four

microcontrollers (μC) AT90CAN128. The CAN controllers on the μC s interact with physical bus via device PCA82C250. The μC throughput achieves 16 MIPS at 16MHz, and the transfer rate of CAN bus is selected at 500kbts/s. By experiment, task *PWM_Calculation* consumes approximately 180 μs to process, the transmission time of *Sensor frame* is approximate 230 μs and the time drift of clock synchronization is 70 μs . Applying formulas (1) and (2), the time slot interval is determined by 500 μs (0.5 ms).

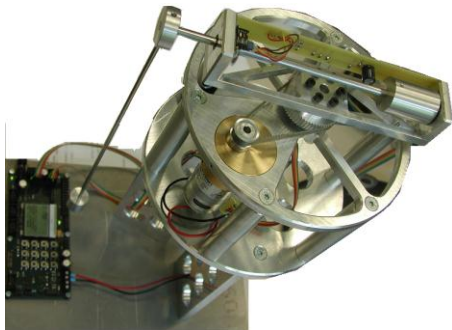


Fig -8: IRP model

The operation system in the μC s is FreeRTOS which is written by Embedded C [5]. It provides the functions to manage the task's implementation and ensure the higher priority task will preemptive the lower one.

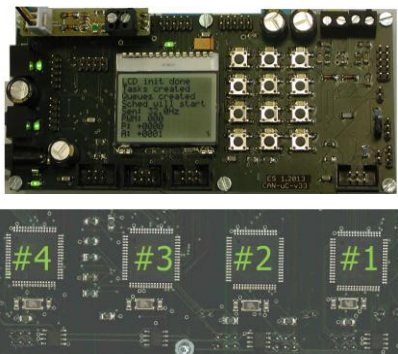


Fig -9: The Control Circuit Board (two sides)

4.2 Results and Discussion

As the IRP model operates at the steady state, the states of TTCAN communication and the hard real-time tasks are measured and shown on Fig -10 and -11. The units of the horizontal and vertical axes are 0.5 milliseconds and 5 voltages, respectively. As the communication channel or tasks are active, the captured signals are non-zero. Because of the repetition of activities, the analysis is carried out in a basic cycle.

On Fig -10, the yellow, green and purple signals represent the states of CAN bus, *PWM_Calculation* task, and the code for preparing data of reference frame (on the third controller), respectively.

The first signal (the yellow one) is CAN bus signal showing the transmission of four message frames in time slot '0', '1', '2', '3' in a basic cycle. The transmission happens as the communication schedule with four message frames in chronological order: MasTim, Sensor, Command, and PWM.

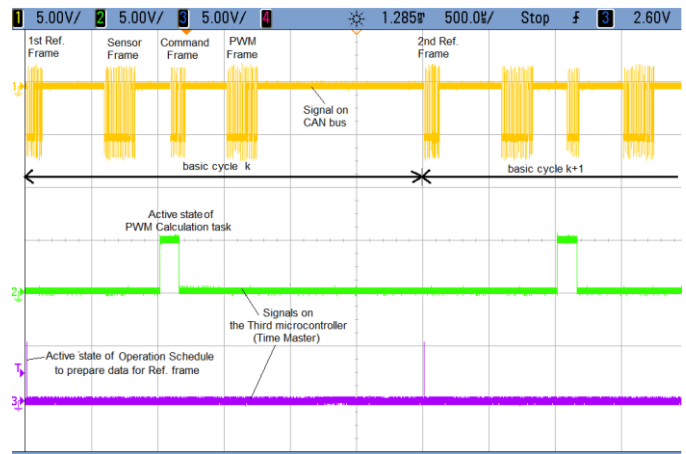


Fig -10: CAN bus signal, states of task *PWM_Calculation* and Part of Operation Schedule

The data for MasTim frame is prepared by a piece of programming codes in *Operation Schedule* task. The state of the codes is captured and shown by the purple signal. It is active periodically at Time Mark '0'. The frame is only sent as its data is ready.

The green signal presents the state of *PWM_Calculation* task that is active in time slot '2' after receiving *Sensor* frame. After completing the calculation, *PWM value* is packed in PWM frame which is sent on time slot '3'. The frame transmission is shown on CAN bus signal.

On CAN bus signal, there are time drift in transmitting two frames *Sensor* and *Command*. The sending points are later than the predefined Time Marks. Due to the delay of clock synchronization, the starting points of basic cycles of Time Master (the third controller) and other μC s are different. Time Master is not affected by this phenomenon. Its tasks and frames' transmission are activated at the right Time Mark.

Using the CAN bus signal as a reference, Fig -11 presents states of tasks *Data Preparation* (first μC - green one) and *Upright Controller* (fourth μC - purple one) to check the synchronization. From the figure, the tasks are activated strictly following the operation schedule. At time slot '0', the task *Data_Preparation* is woken up to prepare the current IRP outputs and pack them into Sensor frame. The frame is then transmitted at Time Mark '1' as shown by signal on CAN bus. *Upright_Controller* task is activated at time slot '4' after receiving PWM frame from the third controller.

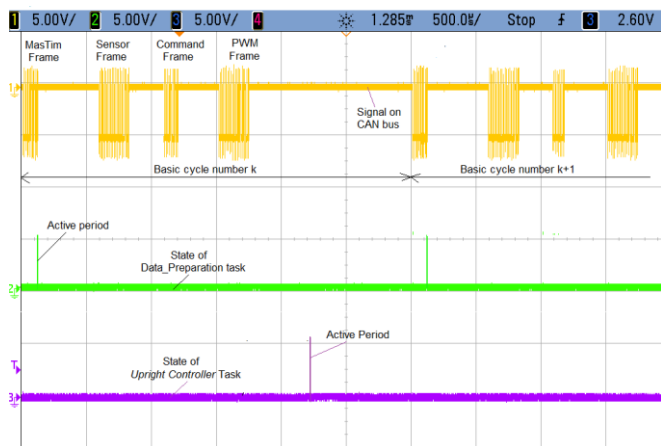


Fig -11: CAN bus signal, states of two tasks Data_Preparation and Upright_Controller

From the figure, the two tasks are woken at middle of the time slot, not at the Time Mark. It is due to the time drift of clock synchronization between Time Master and other uCs.

From the figures above, the critical tasks' operation satisfies the hard real time requirement. They operate synchronously, are active and completed inside the predefined time slots. Fig -12 shows the IRP's operation at steady state. The pendulum can be hold at upright position as long as it is connected to an electrical source.



Fig -12: IRP at steady state

5. CONCLUSIONS

The paper presents a distributed time triggered control for a closed-loop control system. The distributed control meets the real-time requirements and simplifies the structure of the IRP control program. It also can handle both the periodic and aperiodic events deterministically with little latency jitter. A simple method is used to calculate the time slot interval to minimize the basic cycle, thus reduce the effect of the response delay and time drift.

The presented control is deployed for a simple feedback control system with limited environment interaction. The response time to aperiodic events depends on basic cycle. In this experiment, the basic cycle of three milliseconds is able to handle well the operator commands such as stop immediately or adjusting reference input.

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