

A Novel Turbo Codec Encoding and Decoding Mechanism

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Abstract - This paper describes the novel concept of turbo coding and decoding. Basic statistical measures such as an Exclusive OR (XOR) and likelihood were reviewed, and these measures were used to describe the error performance of a decoder using two iterations. We studied the turbo codes working mechanism and how performance is improved when decoders are used in an iterative decoding process using XOR gates. The above technique is been implemented theoretically, first turbo encoder within which RSC encoders working is explained, resultant output is arranged in a sequence. Second turbo decoder using XOR gate is studied and results of second iteration is compared with input sequence and verified

Key Words: Block codes, Convolution codes, Forward error correction, Interleaver, Turbo codes, Turbo decoder, Turbo encoder.

1. INTRODUCTION

Turbo codes offer a way for forward error correction (FEC). In classical turbo encoder [12] the basic form of turbo code generator uses two component codes, separated by an interleaver. Message bits are read into an interleaver by row and then simultaneously read out by rows and by columns into two separate encoders that use either block coding [7] or convolution coding [4] [3]. One encoder is driven by the row message bits and the other by the column message bits from the interleaver, so that an entirely different bit sequence is applied to each encoder, but both encoders are sending the same message bits. Since the row output of the interleaver is the original data stream, one encoder has an input which is the original message bit sequence and the other encoder input is an interleaved version of the message bits. The outputs of the two encoders are combined by multiplying the 2-bit sequences (modulo two). Alternatively, the two outputs can be added and sent sequentially.

Turbo codes based on Convolutional codes are usually known as CTC (Convolutional turbo codes) [1] [8] [9] and those based on block codes as BTC (block turbo codes) [7]. Whereas in turbo codec encoder mechanism the encoder uses a stream-driven implementation and feeds the incoming information bits through to the output. In addition, it uses two recursive convolution encoders (RSC) [5] with one interleaver [4] [3]. The information bits are encoded using encoder-1, a recursive convolution encoder

are parity-1 bits. The interleaved bits are encoded by encoder-2 are parity-2 bits. The encoded bit streams with information bits are arranged in a systematic way. This systematic bit combination is transmitted at the receiver end.

Interleaver is the process of reordering a binary sequence in a systematic way. Convolutional codes are designed to combat random independent errors. However, errors typically come in bursts rather than randomly distributed. Interleaving can be used to disperse the burst errors, making them easier to correct.

NOTE: The turbo codec encoder and decoder interleavers should be identical. The interleaver order used here is as per pseudo random permutations of albert predetermined fashion.

SOVA and log-MAP turbo decoding algorithms are the two prime candidates for decoding turbo codes [11]. In classical turbo decoder [12], at the receiver, the incoming symbol stream is sampled to create a soft input to the two decoders. A soft decoder creates a digital word from each sample using an ADC so that information about the magnitude as well as the state of the received symbol is retained. Recovered bits are given a weighting in the decoding process according to the confidence level from the sampling process. The process is known as soft input soft output (SISO) decoding [11] [12]. In novel turbo codec decoder mechanism we use two decoders and one .. At received data stream, the information bits and parity-1 bits are fed into decoder 1. The interleaved information bits and parity-2 bits are fed into decoder-2. The result of decoder-1 and decoder-2 is iterated two times to get final result.

2. NOVEL TURBO CODEC MECHANISM

We will now illustrate the implementation of a turbo coder using a very simple example. We use the 4 state encoder with block length of only 8 bits.

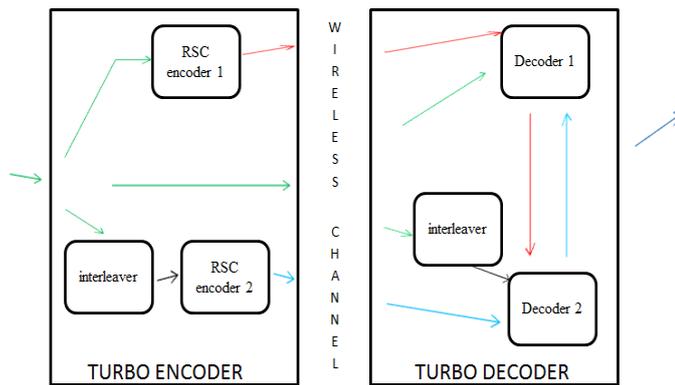


Fig - 1: Block diagram of turbo codec mechanism

2.1. Turbo Encoder:

We are considering the information bits as $\{x_0, x_1, x_2, x_3, x_4, x_5, x_6, x_7\} = \{1, 1, 0, 0, 1, 0, 1, 0\}$ These information bits are passed through turbo encoder and generate output at 1:3 ratios [10] [6].

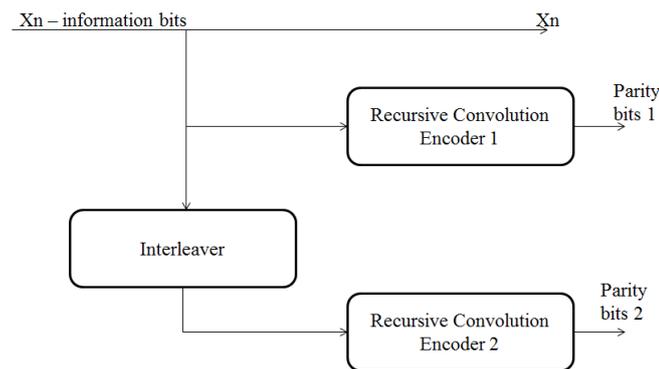


Fig - 2: Block diagram of turbo Encoder

2.1.1. Recursive Convolutional Encoder:

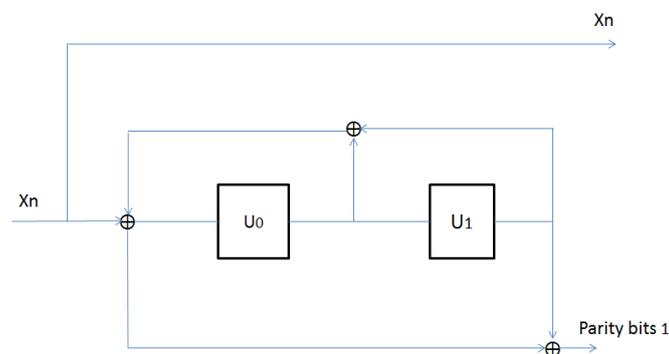


Fig - 3: Recursive convolution encoder internal schematic, State = U1U0

Here U1 and U0 are two flip flops with characteristic equation

$$Q(t+1) \leq Q(t) \tag{1}$$

$$= D \tag{2}$$

Now

U₁ :-

$$Q_1(t+1) = Q_0(t) \tag{3}$$

U₀ :-

$$Q_0(t+1) = X_n \oplus Q_0(t) \oplus Q_1(t) \tag{4}$$

$$= Q_1(t) \oplus Q_1(t) \oplus Q_0(t) \oplus X_n \tag{5}$$

$$= Q_0(t) \oplus X_n \tag{6}$$

Output:-

$$P(x) = Q_0(t) \oplus X_n \tag{7}$$

RSC Encoder 1:

This encoder takes the information bits X_n and generates parity bits which are then passed to decoder 1 in turbo decoder block. The generation of parity bits follows a process using fig 3 where flip-flops with equations (1) to (7) are arranged. Tabulating these values as per equations and collecting the parity bits of recursive convolutional encoder 1 is the same process for the recursive convolutional encoder 2, where encoder1 and encoder 2 parity bits are collected. Let us first see the working of RSC encoder 1 and tabulate the values as follows,

Table -1: RSC Encoder 1 Output

Input X _n	Present state		Next state		Output P ₁ (x) Q ₀ (t) XOR X _n
	Q ₁ (t)	Q ₀ (t)	Q ₁ (t+1)	Q ₀ (t+1)	
1	0	0	0	1	1
1	0	1	1	0	0
0	1	0	0	1	0
0	0	1	1	1	1
1	1	1	1	1	0
0	1	1	1	0	1
1	1	0	0	0	1
0	0	0	0	0	0

Output of RSC encoder 1 is 10010110

Information bits with RSC encoder 1 output combined and the resultant path through the trellis [2] is shown below.

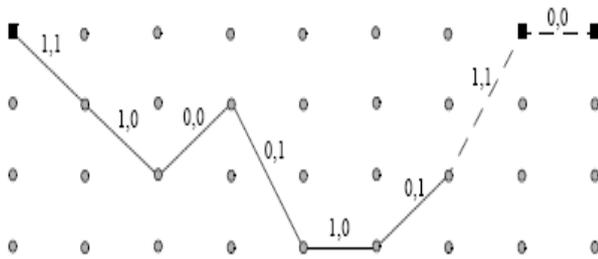


Fig - 4: Trellis path of RSC Encoder 1

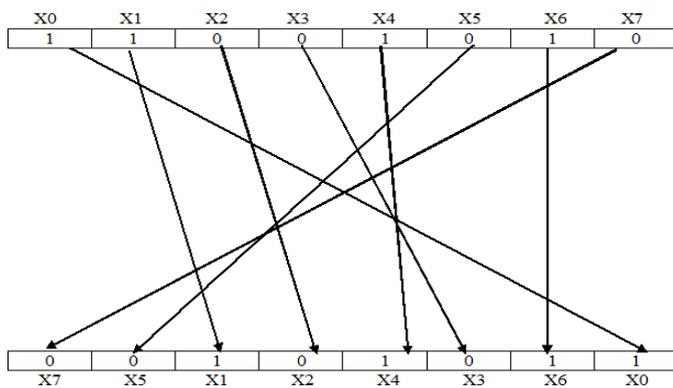


Fig - 5: Pseudo random interleaver (predefined pattern)

The interleaver pattern is predefined and identical to the both sides of turbo codec. The sequence of the information bits is shuffled using this interleaver pattern. This interleaved information bits is passed through the recursive convolutional encoder 2 to generate parity bits. This follows the same process by using fig 3 and equation (1) to (7). We tabulate the bit values using XOR operator. These parity bits are arranged systematically with the information bits and parity bits of encoder 1 and form a 24 bit sequence.

Table -2: RSC Encoder 2 Output

Interleaved Input X_{nk}	Present state		Next state		Output P1 (x) $Q_0(t) \text{ XOR } X_n$
	$Q_1(t)$	$Q_0(t)$	$Q_1(t+1)$	$Q_0(t+1)$	
0	0	0	0	0	0
0	0	0	0	0	0
1	0	0	0	1	1
0	0	1	1	1	1
1	1	1	1	1	0
0	1	1	1	0	1
1	1	0	0	0	1
1	0	0	0	1	1

RSC encoder 2 output is 00110111 interleaved information bits with RSC encoder 2 output resulting a path through trellis [2].

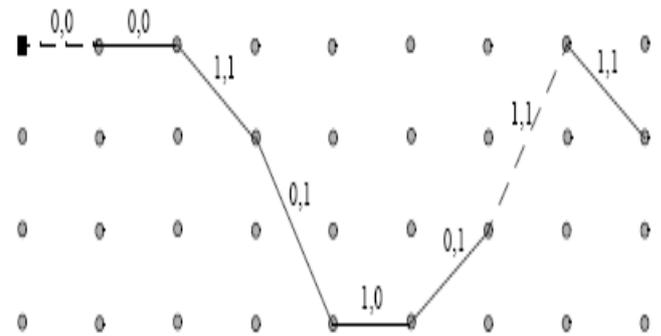


Fig - 6: Trellis path of RSC Encoder 2

2.1.2. Arrangement of information bits and parity bits:

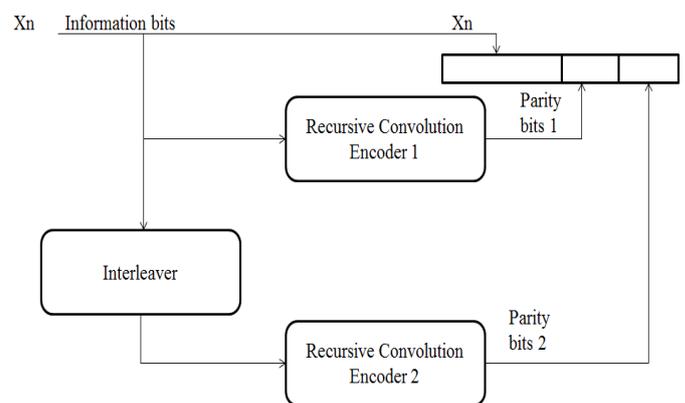


Fig - 7: Arrangement of information bits, parity bits 1 (encoder 1 output bits) and parity bits 2 (encoder 2 output bits)

Information bits	Parity 1 bits	Parity bits 2
11001010	10010110	0011011

This 24 bits data is transmitted through channel and received at the receiving antenna and the received code word will be 110010101001011000110111

2.2 Turbo Decoder:

We now feed the block of data through turbo decoder. Turbo decoder consists of two decoders and one interleaver. The interleaver order used here is as per pseudo random permutations of albert predetermined fashion. Identical interleavers should be used in turbo codec encoder and decoder.

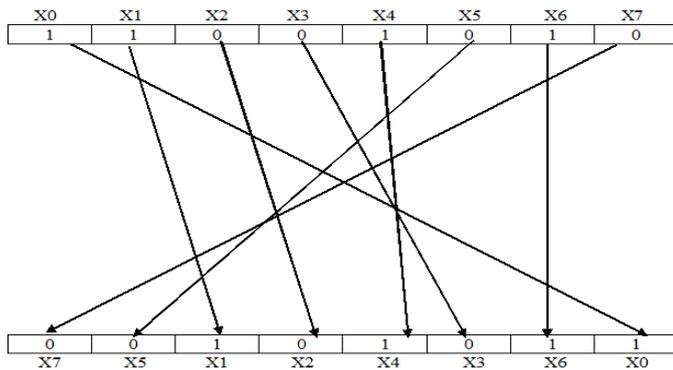


Fig - 8: Interleaver at turbo decoder block

Interleaved output = 00101011

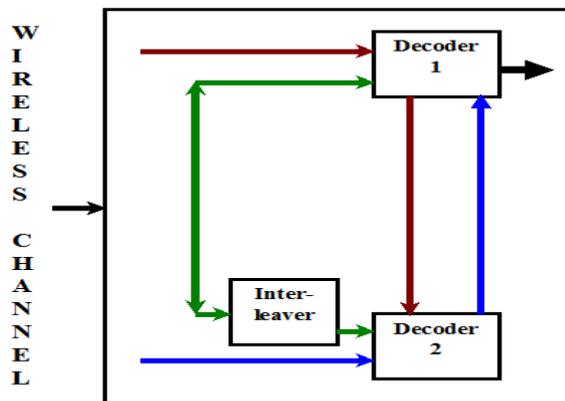


Fig - 9: Block diagram of Turbo decoder

Information bits = 11001010
 Parity bits 1 or Encoder 1 bits = 10010110
 Parity bits 2 or Encoder 2 bits = 00110111
 Interleaver = 00101011

2.2.1. Working of Decoder using XOR gate:-

At received data stream, the information bits and parity 1 bits are fed into decoder 1. The interleaved information bits and parity 2 bits are fed into decoder 2. The result of decoder 1 and decoder 2 is iterated two times to get final results.

Note:- “⊕” is XOR gate operator.

Decoder 1:-

Information bits ⊕ parity 1 bits
 11001010
 ⊕ 10010110

 01011100 -- Decoder 1 output

Decoder 2:-

Interleaved information bits ⊕ parity 2 bits
 00101011
 ⊕ 00110111

 00011100 -- Decoder 2 output

First Iteration:-

Information bits ⊕ decoder 1 ⊕ decoder 2
 11001010
 ⊕ 01011100
 ⊕ 00011100

 10001010 -- First iteration output

Second iteration:-

First iterated output ⊕ decoder 1 ⊕ decoder 2
 10001010
 ⊕ 01011100
 ⊕ 00011100

 11001010 -- Second iterated output = original input bits.

Hence we got the input information bits sequence at decoder in second iteration itself, which reduces latency and complexity.

3. CONCLUSION

This paper has proposed, for the first time, a novel turbo codec method of turbo codes which reduces the complexity of turbo decoder working by using a simple XOR gate. In particular, a new approach with two iteration and getting the original information signal. In earlier decoder algorithms SOVA and log-MAP turbo decoding algorithms were used for decoding turbo codes. On comparison, it is noted that proposed novel turbo codec mechanism reduces complexity. Turbo encoder detailed working mechanism is given in this paper. The interleavers were predetermined valued interleavers which is identical at both side of turbo codec mechanism.

It is under the current investigation to extend the proposed methodology to the satellites using FPGAs. The above technique can be implemented using C-code simulation and VHDL for FPGAs.

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BIOGRAPHIES

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